



Chapter 1: Introduction

These slides are a modified version of the slides provided with the book:

Database System Concepts, 6th Ed.

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Outline

- The Need for Databases
- Data Models
- Relational Databases
- Database Design
- Storage Manager
- Query Processing
- Transaction Manager



Database Management System (DBMS)

- DBMS contains information about a particular enterprise
 - Collection of interrelated data
 - Set of programs to access the data
 - An environment that is both *convenient* and *efficient* to use
- Database Applications:
 - Banking: transactions
 - Airlines: reservations, schedules
 - Universities: registration, grades
 - Sales: customers, products, purchases
 - Online retailers: order tracking, customized recommendations
 - Manufacturing: production, inventory, orders, supply chain
 - Human resources: employee records, salaries, tax deductions
- Databases can be very large.
- Databases touch all aspects of our lives



University Database Example

- Application program examples
 - Add new students, instructors, and courses
 - Register students for courses, and generate class rosters
 - Assign grades to students, compute grade point averages (GPA) and generate transcripts
- In the early days, database applications were built directly on top of file systems



Drawbacks of using file systems to store data

- Data redundancy and inconsistency
 - Multiple file formats, duplication of information in different files
- Difficulty in accessing data
 - Need to write a new program to carry out each new task
- Data isolation
 - Multiple files and formats
- Integrity problems
 - Integrity constraints (e.g., account balance > 0) become “buried” in program code rather than being stated explicitly
 - Hard to add new constraints or change existing ones



Drawbacks of using file systems to store data (Cont.)

- Atomicity of updates
 - Failures may leave database in an inconsistent state with partial updates carried out
 - Example: Transfer of funds from one account to another should either complete or not happen at all
- Concurrent access by multiple users
 - Concurrent access needed for performance
 - Uncontrolled concurrent accesses can lead to inconsistencies
 - ▶ Example: Two people reading a balance (say 100) and updating it by withdrawing money (say 50 each) at the same time
- Security problems
 - Hard to provide user access to some, but not all, data

Database systems offer solutions to all the above problems



Levels of Abstraction

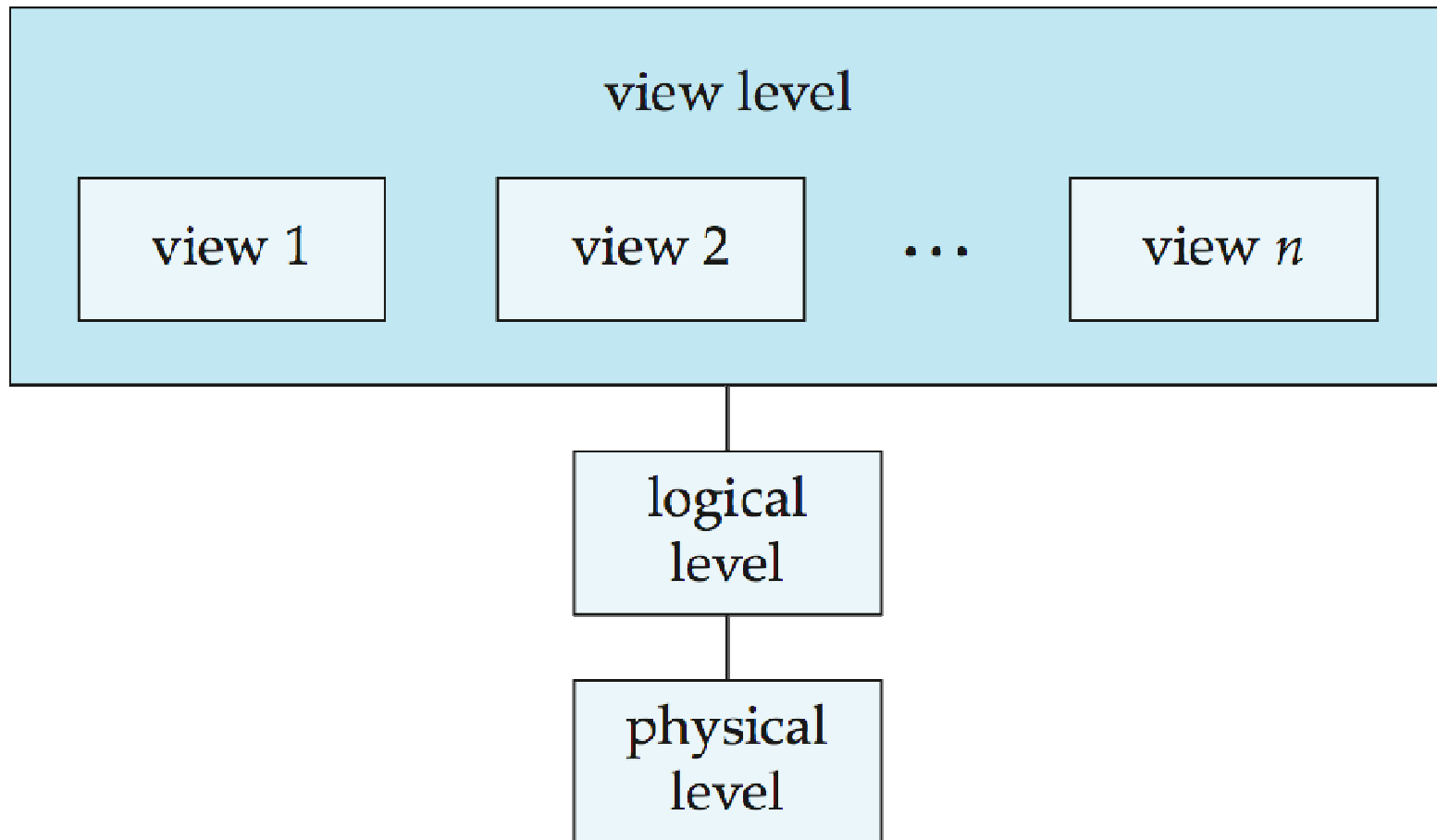
- **Physical level:** describes how data is physically stored in files.
- **Logical level:** describes data stored in database, and the relationships among the data.
- **View level:** application programs hide details of data types. Views can also hide information (such as an employee's salary) for security purposes.

- **Data Independence (transparency)**
 - In general, the interfaces between the various levels and components should be well defined so that changes in some parts do not seriously influence others.
 - **Physical** data independence (before referred to as **data isolation**) – the ability to modify the physical schema without changing the logical schema
 - **Logical** data independence – the ability to modify the logical schema without changing applications and user interfaces, and vice versa



View of Data

An architecture for a database system





Data Models

- A collection of tools for describing
 - Data
 - Data relationships
 - Data semantics
 - Data constraints
- Relational model
- Entity-Relationship data model (mainly for database design)
- Object-based data models (Object-oriented and Object-relational)
- Semistructured data model (XML)
- Other older models:
 - Network model
 - Hierarchical model



Relational Model

- All the data is stored in various tables.
- Example of tabular data in the relational model

Columns

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

Rows

(a) The *instructor* table



Instances and Schemas

- **Relation Schema** – the list of attributes and their domains

```
type instructor = record
```

```
    ID : string;
```

```
    name : string;
```

```
    dept_name : string;
```

```
    salary : integer;
```

```
end;
```

- Analogous to the type of a variable
- **Relation Instance** – the actual content of the database at a particular point in time (a set of tuples)
 - Analogous to the value of a variable
- **DB Schema** – the set of relation schemas
- **DB Instance** – the set of relation instances



A Sample Relational Database

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
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33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

<i>dept_name</i>	<i>building</i>	<i>budget</i>
Comp. Sci.	Taylor	100000
Biology	Watson	90000
Elec. Eng.	Taylor	85000
Music	Packard	80000
Finance	Painter	120000
History	Painter	50000
Physics	Watson	70000

(b) The *department* table



Data Definition Language (DDL)

- Specification notation for defining the database schema

Example: **create table** *instructor* (
 ID **char**(5),
 name **varchar**(20),
 dept_name **varchar**(20),
 salary **numeric**(8,2))

- DDL compiler generates a set of table templates stored in a ***data dictionary***
- Data dictionary contains metadata (i.e., data about data)
 - Database schema
 - Integrity constraints
 - ▶ Primary key (ID uniquely identifies instructors)
 - Authorization
 - ▶ Who can access what



Data Manipulation Language (DML)

- Language for accessing and manipulating the data organized by the appropriate data model
 - DML also known as query language
- Two classes of languages
 - **Pure** – used for proving properties about computational power and for optimization
 - ▶ Relational Algebra
 - ▶ Tuple relational calculus
 - ▶ Domain relational calculus
 - **Commercial** – used in commercial systems
 - ▶ SQL is the most widely used commercial language



SQL

- The most widely used commercial language
- SQL is NOT a Turing machine equivalent language
- To be able to compute complex functions SQL is usually embedded in some general-purpose language
- Application programs generally access databases through one of
 - Language extensions to allow embedded SQL
 - Application program interface (e.g., ODBC/JDBC) which allow SQL queries to be sent to a database



Database Design

The process of designing the general structure of the database:

- Logical Design – Deciding on the database schema.
Database design requires that we find a “good” collection of relation schemas.
 - Business decision – What attributes should we record in the database?
 - Computer Science decision – What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
- Physical Design – Deciding on the physical layout of the database



Database Design (Cont.)

- Is there any problem with this relation?

<i>ID</i>	<i>name</i>	<i>salary</i>	<i>dept_name</i>	<i>building</i>	<i>budget</i>
22222	Einstein	95000	Physics	Watson	70000
12121	Wu	90000	Finance	Painter	120000
32343	El Said	60000	History	Painter	50000
45565	Katz	75000	Comp. Sci.	Taylor	100000
98345	Kim	80000	Elec. Eng.	Taylor	85000
76766	Crick	72000	Biology	Watson	90000
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33456	Gold	87000	Physics	Watson	70000
76543	Singh	80000	Finance	Painter	120000



Design Approaches

- Need to come up with a methodology to ensure that each of the relations in the database is “good”
- Two ways of doing so:
 - Entity Relationship Model (Chapter 7)
 - ▶ Models an enterprise as a collection of *entities* and *relationships*
 - ▶ Represented diagrammatically by an *entity-relationship diagram*:
 - Normalization Theory (Chapter 8)
 - ▶ Formalize what designs are bad, and test for them



DBMS Components

- Storage manager
- Query processing
- Transaction manager



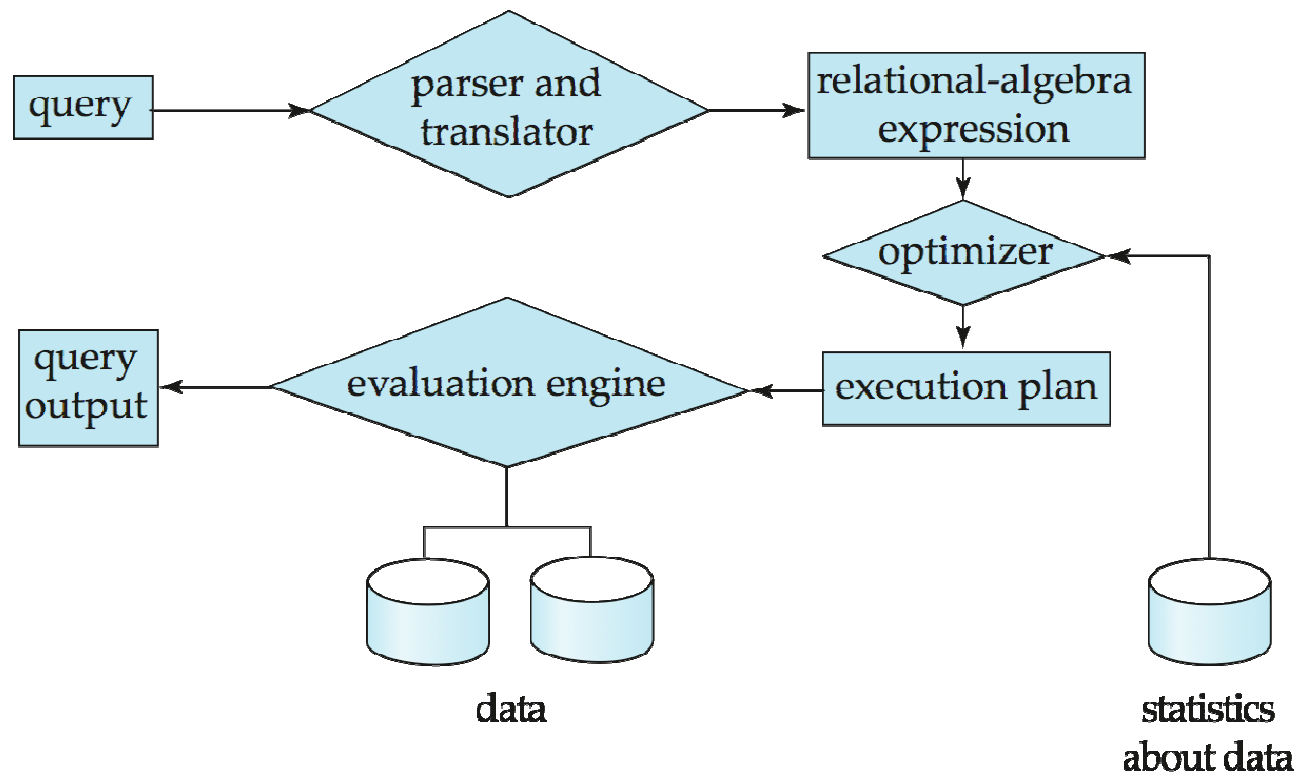
Storage Management

- **Storage manager** is a program module that provides the interface between the low-level data stored in the database and the application programs and queries submitted to the system.
- The storage manager is responsible to the following tasks:
 - Interaction with the OS file manager
 - Efficient storing, retrieving and updating of data
- Issues:
 - Storage access
 - File organization
 - Indexing and hashing



Query Processing

1. Parsing and translation
2. Optimization
3. Evaluation





Query Processing (Cont.)

- Alternative ways of evaluating a given query
 - Equivalent expressions
 - Different algorithms for each operation
- Cost difference between a good and a bad way of evaluating a query can be enormous
- Need to estimate the cost of operations
 - Depends critically on statistical information about relations which the database must maintain
 - Need to estimate statistics for intermediate results to compute cost of complex expressions

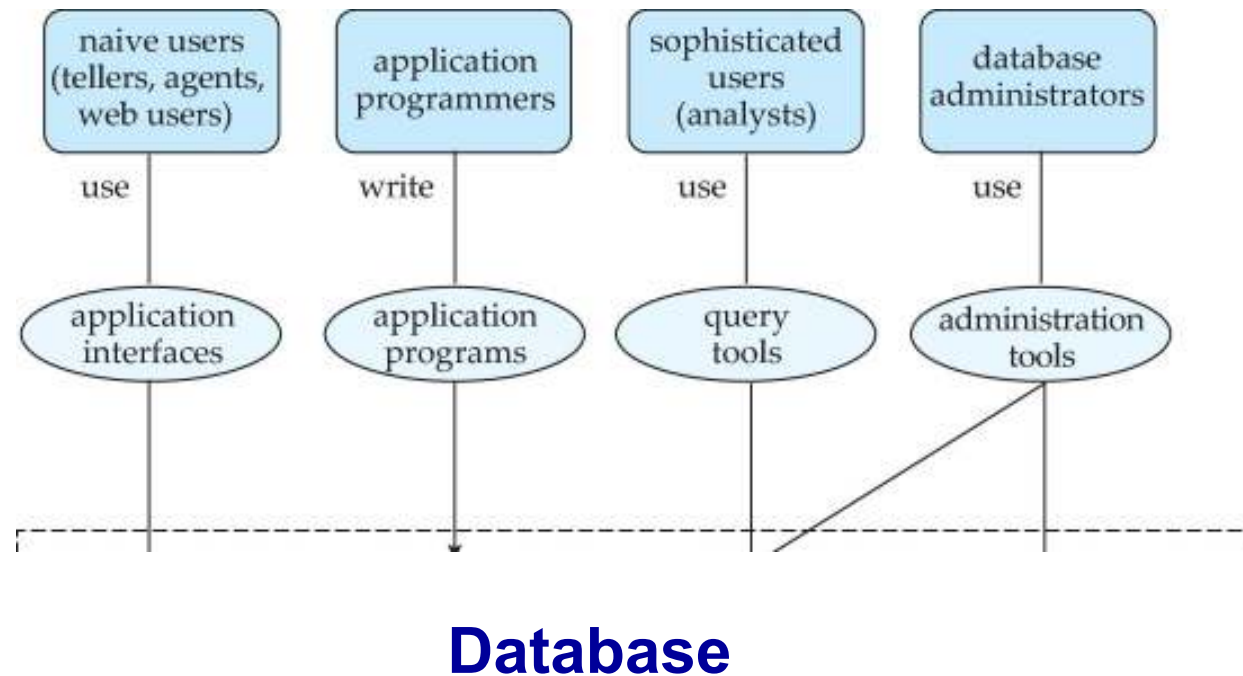


Transaction Management

- What if the system fails?
- What if more than one user is concurrently updating the same data?
- A **transaction** is a collection of operations that performs a single logical function in a database application
- **Transaction-management component** ensures that the database remains in a consistent (correct) state despite system failures (e.g., power failures and operating system crashes) and transaction failures.
- **Concurrency-control manager** controls the interaction among the concurrent transactions, to ensure the consistency of the database.

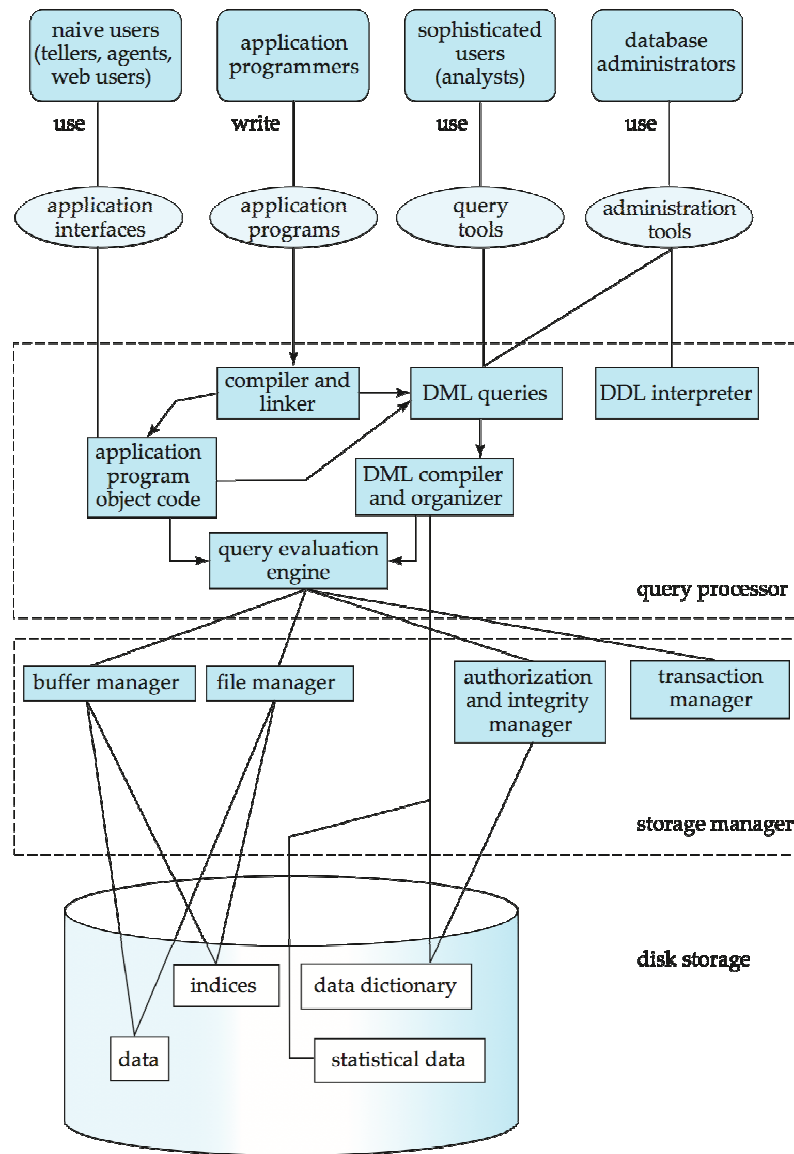


Database Users and Administrators





Database System Internals





Database Architecture

The architecture of a database systems is greatly influenced by the underlying computer system on which the database is running:

- Centralized
- Client-server
- Parallel (multi-processor)
- Distributed



History of Database Systems

- 1950s and early 1960s:
 - Data processing using magnetic tapes for storage
 - ▶ Tapes provided only sequential access
 - Punched cards for input
- Late 1960s and 1970s:
 - Hard disks allowed direct access to data
 - Network and hierarchical data models in widespread use
 - Ted Codd defines the relational data model
 - ▶ Would win the ACM Turing Award for this work
 - ▶ IBM Research begins System R prototype
 - ▶ UC Berkeley begins Ingres prototype
 - High-performance (for the era) transaction processing



History (cont.)

- 1980s:
 - Research relational prototypes evolve into commercial systems
 - ▶ SQL becomes industrial standard
 - Parallel and distributed database systems
 - Object-oriented database systems
- 1990s:
 - Large decision support and data-mining applications
 - Large multi-terabyte data warehouses
 - Emergence of Web commerce
- Early 2000s:
 - XML and XQuery standards
 - Automated database administration
- Later 2000s:
 - Giant data storage systems
 - ▶ Google BigTable, Yahoo PNuts, Amazon, ..



End of Chapter 1