

Model Checking Coalitional Games with Priced-Resource Agents

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At a glance

- We presented the logic PRB-ATL [CILC 11 & GAMES 11]
- PRB-ATL is inspired to existing extensions of ATL
 - ▶ To deal with bounded resources scenarios in multi-agent systems
- We studied the model checking for PRB-ATL
 - ▶ It is in EXPTIME (upper bound) [LAMAS/M4M 11]
 - ▶ It is EXPTIME-hard (lower bound) [GAMES 12]
- We studied the optimal coalition problem [LAMAS/M4M 11]

- 1 Introduction to Multi-Agent Systems (MAS) - **ATL**
 - Multi-Agent Systems and resource constraints - **RB-ATL**
- 2 Our proposal: the logic *Priced RB-ATL* - **PRB-ATL**
 - Model checking
 - Optimization problem
- 3 Conclusions

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Multi-Agent Systems (MAS)

- Several agents
- Intelligent (take decision)
- Independent
- Global state (union of single states)
- Move choices
- Next state

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COALITION - modeling collective behaviors/strategies

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Logical Formalisms

Coalition Logic (CL) and Alternating-time Temporal Logic (ATL)

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Logical Formalisms

Coalition Logic (CL) and Alternating-time Temporal Logic (ATL)

Theorem (Goranko, TARK 2001)

CL can be embedded into ATL

ATL: syntax and semantics

Formulae of ATL are given by the grammar:

$$\varphi ::= p \mid \neg\varphi \mid \varphi \wedge \varphi \mid \langle\langle A \rangle\rangle \bigcirc \varphi \mid \langle\langle A \rangle\rangle \varphi \mathcal{U} \varphi \mid \langle\langle A \rangle\rangle \square \varphi$$

Formulae of ATL predicate about abilities of coalitions of agents

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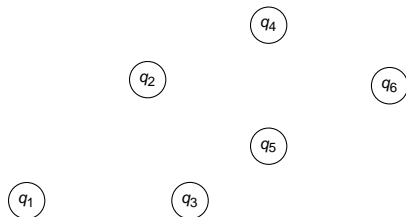
Formulae of ATL predicate about abilities of coalitions of agents

Formulae of ATL are evaluated wrt:

- a **game structure** (or **game arena**) G
- a **location** q of G

The arena of ATL

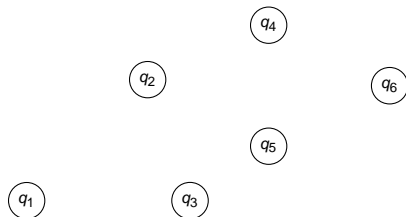
A **game structure** G is a state transition graph:



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- in each location, each agent can choose among a non-empty set of **actions**
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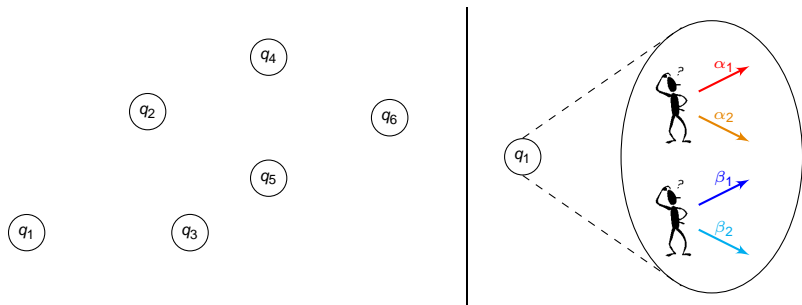
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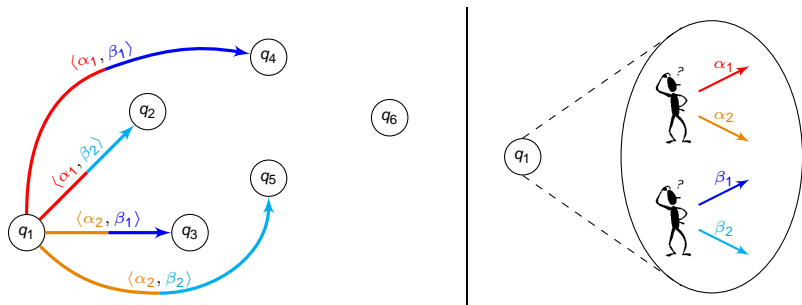
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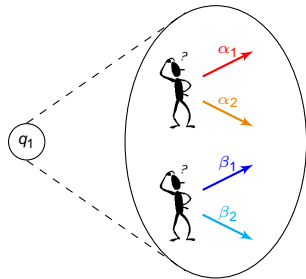
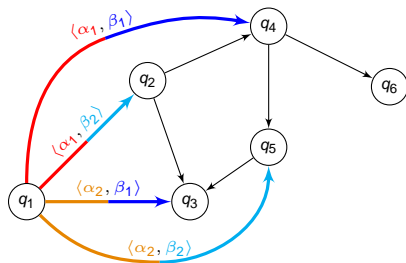
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Becoming friendly with ATL

Collective strategy to guarantee p holds

$\langle\langle A \rangle\rangle \bigcirc p$ next

$\langle\langle A \rangle\rangle \square p$ always

$\langle\langle A \rangle\rangle p \mathcal{U} q$ until q

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regardless of actions performed by other agents (opponent)

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Addition of bounds on resources to ATL



o' munn' è piccirill!!!



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Quant'è piccolo il mondo!

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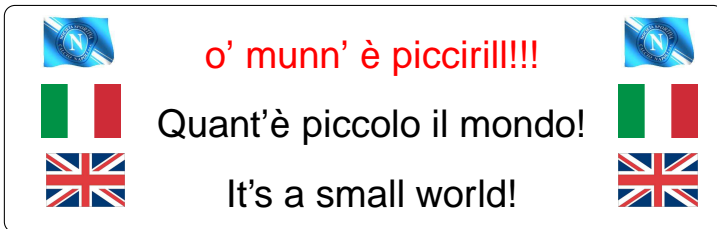
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Quant'è piccolo il mondo!

It's a small world!



Addition of bounds on resources to ATL



Resources are bounded

Extensions of ATL with bounds on resources

$$\langle\langle A^\eta \rangle\rangle \Box p$$

Endowment: $\eta : A \rightarrow \mathbb{N}^r$

The literature about Resource Bounded ATL (RB-ATL)

RB-ATL [Alechina, Logan, Nga, Rakib, AAMAS 2010]

Theorem: Model checking RB-ATL is decidable in $O(|\varphi|^{2 \cdot r + 1} \times |G|)$
No lower bound

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RAL [Bulling, Farwer, ECAI 2010]

If actions may produce resources,
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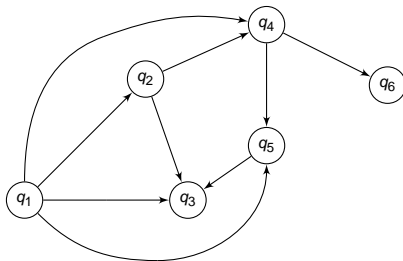
Formulae of RB-ATL predicate about abilities of coalitions whose agents are equipped with a finite endowment of resources

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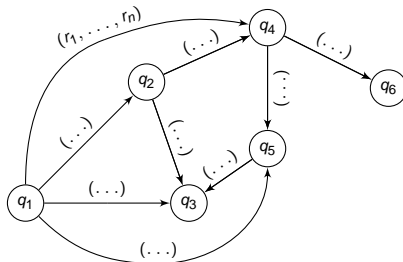
A **resource-bounded game structure G** is a **weighted** state transition graph:



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Becoming friendly with RB-ATL

$$G, q \Vdash \langle\langle A^\eta \rangle\rangle \bigcirc \langle\langle A^{\eta'} \rangle\rangle \Box p$$

team A , equipped with endowment η , can force the next state to be s.t. the team A itself can guarantee that p always holds equipped with the new endowment η'

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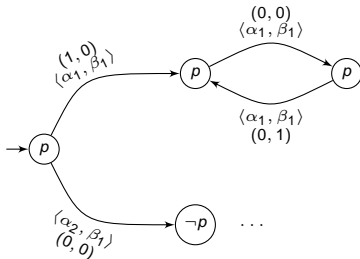
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An anomalous behavior

2 agents: ag_1 and ag_2
1 resource type: r_1

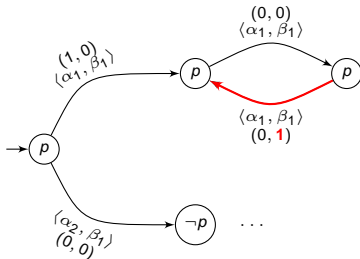
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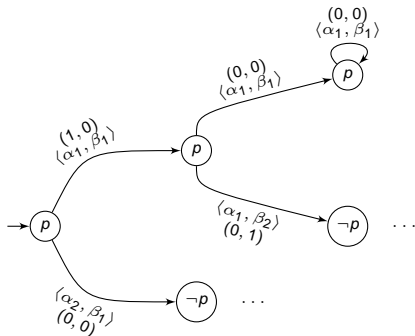


opponent consumes an infinite amount of resources

Another anomalous behavior

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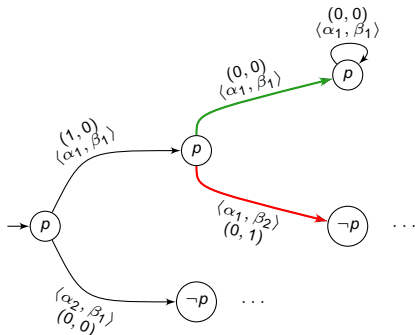
$$G, q_0 \Vdash \langle\langle ag_1^? \rangle\rangle \Box p$$



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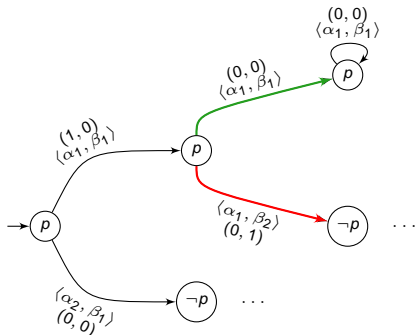


false if $r_1 > 1$
true if $r_1 = 1$

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opponent's moves should be constrained

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Weaknesses of previous approaches

- NO history (resources)

- ▶ $G, q \Vdash \langle\langle A^\eta \rangle\rangle \circ \langle\langle A^{\eta'} \rangle\rangle \Box p$ η and η' are independent

- opponent does NOT consume

- ▶ opponent has no bounds on resources
- ▶ consumption by opponent does not matter

What we want

- opponent's actions constrained
- consumption/production tracked
- a significant present-day issue \Rightarrow procurement of resources
 - ▶ **limited amount** on the market (or in nature)
 - ▶ **acquisition cost** depending on current availability

How we get it

Key notion \Rightarrow **global availability of resources on the market**

- a semantic component (part of the arena)
- evolves depending on agents' actions (also opponent)
- affects the choice of the actions (also opponent)

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Auxiliary notion \Rightarrow **price of resources**

- agents equipped with money instead of resources
- money for getting resources
- price of resources function of several components (take into account the history of the system)

Money vs. resources - our proposal

Money

- inside the **formula**
- assigned to **agents**
- **private**: any agent has his own amount of money
- **unknown**
- availability checked for **proponent's agents only**

Resources

- part of the **model**
- represent the **market** (**nature**)
- **public**: agents draw on resources from a shared pool
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Money is a *meta-resource*

- buy resources
 - ▶ money like resources in previous approaches
- **unit of measurement**

Resource production and decidability

Alechina, Logan, Nga, Rakib

Actions can **only consume** resources

Bulling, Farwer

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Actions may produce resources...
...but **not so much!!!**

- **model checking decidable**
- **several models fit** (memory usage, leasing a car, releasing resources previously acquired)

Syntax and semantics

Formulae of PRB-ATL are given by the grammar:

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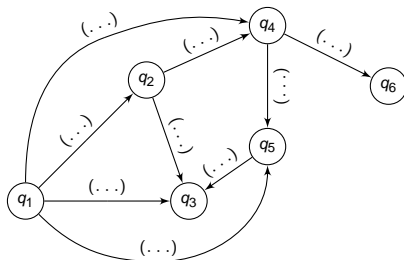
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Priced game structure

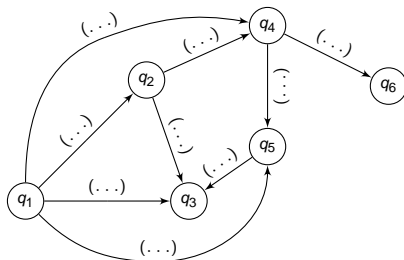
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- resources have a variable **prices**
- transition guards: **also opponent**

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Model checking complexity

Theorem

The model checking problem for PRB-ATL is EXPTIME-complete

- *membership (upper bound):* [LAMAS 2011]
- *hardness (lower bound):* [GAMES 2012]

- Same asymptotic complexity as RB-ATL
- Exponential blow-up compared to ATL

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Reduction from the acceptance problem for
Linearly-Bounded Alternating Turing Machine

Reduction

- Given a LB-ATM \mathcal{A} and an input s_0
- We provide
 - ▶ a priced game structure $G_{\mathcal{A}, s_0}$
 - ▶ an initial location q_0 in $G_{\mathcal{A}, s_0}$
 - ▶ an initial availability of resources \vec{m}_0
 - ▶ a PRB-ATL formula $\varphi_{\mathcal{A}, s_0}$

such that

$$G_{\mathcal{A}, s_0}, q_0, \vec{m}_0 \models \varphi_{\mathcal{A}, s_0} \text{ iff } \mathcal{A} \text{ accepts on input } s_0$$

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Parametric PRB-ATL formulae

- PRB-ATL: $\varphi = \langle\langle A_1^{\$1} \rangle\rangle \diamond (\langle\langle A_2^{\$2} \rangle\rangle \circ p \vee \langle\langle A_3^{\$3} \rangle\rangle q \cup p)$

Definition (Cost of a PRB-ATL formula)

$$f_cost(\varphi) = \$_1(A_1) + \$_2(A_2) + \$_3(A_3)$$

- parametric PRB-ATL: $\varphi_{\vec{x}} = \langle\langle X_1^{\$1} \rangle\rangle \diamond (\langle\langle X_2^{\$2} \rangle\rangle \circ p \vee \langle\langle A_3^{\$3} \rangle\rangle q \cup p)$

The *Optimal Coalition* problem

Definition (Optimal Coalition problem)

To determine minimal-cost coalitions that satisfy a PRB-ATL formula

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Conclusions

A logic for modeling multi-agent systems with bounds on resources

ATL: abilities of coalitions of agents

RB-ATL: abilities of coalitions whose agents are equipped with a finite **endowment of resources**

PRB-ATL: abilities of coalitions whose agents are equipped with an amount of **money**

- **global availability of resources**
- **money - price of resources**

- Resource-bounded extensions of other classical formalisms
 - ▶ e.g., μ -calculus [Della Monica, Lenzi - ICAART 2012]
- Hierarchical, MAS, and resources???
- An attempt towards tractability: the *reachability problem*

Future work

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