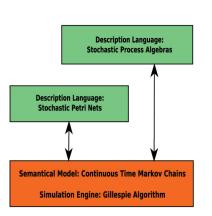
Modeling Biological Systems in Stochastic Concurrent Constraint Programming

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Modeling Biological Systems with Stochastic Process Algebras



Pros

- Simple Language
- Compositionality

Cons

- Hard to encode general information
- Lacking computational extensibility

Outline

- Theory
 - Concurrent Constraint Programming
 - Continuous Time Markov Chains
 - Stochastic CCP
- Bio-Modeling
 - Modeling Biochemical Reactions
 - Modeling Gene Regulatory Networks

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Concurrent Constraint Programming

Constraint Store

- In this process algebra, the main object are constraints, which are formulae over an interpreted first order language (i.e. X = 10, Y > X - 3).
- Constraints can be added to a "pot", called the constraint store, but can never be removed.

Agents

Agents can perform two basic operations on this store:

- Add a constraint (tell ask)
- Ask if a certain relation is entailed by the current configuration (ask instruction)

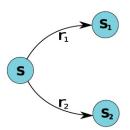
Syntax of CCP

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\begin{aligned} &\textit{Program} = \textit{Decl.A} \\ &\textit{D} = \varepsilon \mid \textit{Decl.Decl} \mid \textit{p(x)} : -\textit{A} \\ &\textit{A} = \quad \mathbf{0} \\ &\mid \quad \text{tell(c).A} \\ &\mid \quad \text{ask(c_1).A}_1 + \text{ask(c_2).A}_2 \\ &\mid \quad A_1 \mid\mid A_2 \mid \exists_{\textit{x}} \textit{A} \mid \textit{p(x)} \end{aligned}
```



Continuous Time Markov Chains

A **Continuous Time Markov Chain** (CTMC) is a direct graph with edges labeled by a real number, called the **rate of the transition** (representing the **speed** or the **frequency** at which the transition occurs).



- The time spent in a state is given by an exponentially distributed random variable, with rate given by the sum of outgoing transitions from the actual node (r₁ + r₂).



Syntax of sCCP

Syntax of Stochastic CCP

$$Program = D.A$$

$$D = \varepsilon \mid D.D \mid p(\mathbf{x}) : -A$$

$$\pi = \operatorname{tell}_{\lambda}(c) \mid \operatorname{ask}_{\lambda}(c)$$

$$M = \pi.A \mid \pi.A.p(\mathbf{y}) \mid M + M$$

$$A = \mathbf{0} \mid \operatorname{tell}_{\infty}(c).A \mid \exists_{\mathbf{x}} A \mid M \mid (A \parallel A)$$

Stochastic Rates

Each basic instruction (tell, ask, procedure call) has a rate attached to it. Rates are functions from the constraint store \mathcal{C} to positive reals: $\lambda: \mathcal{C} \longrightarrow \mathbb{R}^+$.



sCCP soup

Operational Semantics

- There are two transition relations, one instantaneous (finite and confluent) and one stochastic.
- Traces are sequences of events with variable time delays among them.

Implementation

- We have an interpreter written in Prolog, using the CLP engine of SICStus to manage the constraint store.
- Efficiency issues.

Stream Variables

- Quantities varying over time can be represented in sCCP as unbounded lists.
- Hereafter: special meaning of X = X + 1.

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General Principles

 $\mbox{Measurable Entities} \quad \leftrightarrow \quad \mbox{Stream Variables}$

Logical Entities

→ Processes (Control Variables)

Interactions \leftrightarrow Processes

Biochemical Arrows to sCCP processes

$$R_{1} + \ldots + R_{n} \rightarrow_{k} P_{1} + \ldots + P_{m}$$

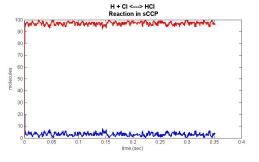
$$= \frac{k_{1}}{k_{N}} (k_{1}, R_{1}, \ldots, R_{n}), (\bigcap_{i=1}^{n} (R_{i} > 0)) \cdot (\prod_{i=1}^{m} \text{tell}_{\infty}(P_{i} = P_{i} + 1)) \cdot \prod_{i=1}^{m} \text{tell}_{\infty}(P_{i} = P_{i} + 1)) \cdot \prod_{i=1}^{m} \text{tell}_{\infty}(P_{i} = P_{i} + 1) \cdot \prod_{i=1}^{m} \text{tell}_{\infty}(P_{i} = P_{i} + 1)) \cdot \prod_{i=1}^{m} \text{tell}_{\infty}(P_{i} = P_{i} + 1) \cdot \prod_{i=1}^{m} \text{tell}_{\infty}(P_{i} = P_$$

$$\text{where} \quad r_{MA}(k,X_1,\ldots,X_n) = k \cdot X_1 \cdot \cdots \cdot X_n; \quad r_{MM}(K,V_0,S) = \frac{V_0 \, S}{S + K}; \quad r_{Hill}(k,V_0,h,S) = \frac{V_0 \, S^h}{S^h + K^h}$$

A simple reaction: $H + CI \leftrightharpoons HCI$

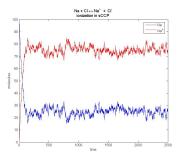
We have two reaction agents. The reagents and the products are stream variables of the constraint store (put down in the environment). *Independent on the number of molecules*.

 $\mathrm{reaction}(100,[H,CL],[HCL]) \parallel \mathrm{reaction}(10,[HCL],[H,CL])$



Another reaction: Na + Cl \rightleftharpoons Na⁺ + Cl⁻

 $\mathrm{reaction}(100, [\textit{NA}, \textit{CL}], [\textit{NA}+, \textit{CL}-]) \parallel \mathrm{reaction}(10, [\textit{NA}+, \textit{CL}-], [\textit{NA}, \textit{CL}])$



Enzymatic reaction

$$S + E \rightleftharpoons_{k_{-1}}^{k_1} ES \rightarrow_{k_2} P + E$$

Mass Action Kinetics

enz_reaction(k_1 , k_{-1} , k_2 , S, E, ES, P):reaction(k_1 , [S, E], [ES]) \parallel reaction(k_{-1} , [ES], [E, S]) \parallel reaction(k_2 , [ES], [E, P])

Mass Action Equations

$$\begin{array}{l} \frac{d[ES]}{dt} = k_1[S][E] - k_2[ES] - k_{-1}[ES] \\ \frac{d[E]}{dt} = -k_1[S][E] + k_2[ES] + k_{-1}[ES] \\ \frac{d[S]}{dt} = -k_1[S][E] \\ \frac{d[P]}{dt} = k_2[ES] \end{array}$$

Michaelis-Menten Equations

$$\frac{d[P]}{dt} = \frac{V_0 S}{S+K}$$

$$V_0 = k_2 [E_0]$$

$$K = \frac{k_2 + k - 1}{k_4}$$

Michaelis-Menten Kinetics

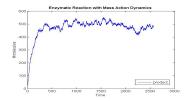
mm_reaction
$$\left(\frac{k_2 + k_{-1}}{k_1}, k_2 \cdot E, S, P\right)$$

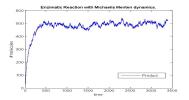
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enz_reaction(k_1 , k_{-1} , k_2 , S, E, ES, P):reaction(k_1 , [S, E], [ES]) \parallel reaction(k_{-1} , [ES], [E, S]) \parallel reaction(k_0 , [ES], [E, P])

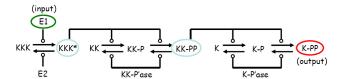




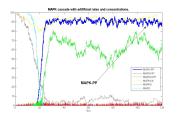
Michaelis-Menten Kinetics

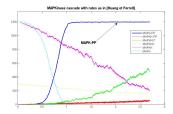
mm_reaction
$$\left(\frac{k_2 + k_{-1}}{k_1}, k_2 \cdot E, S, P\right)$$

MAP-Kinase cascade

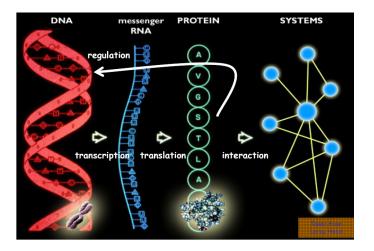


enz_reaction(k_a, k_d, k_r , KKK, E1, KKKE1, KKKS) \parallel enz_reaction(k_a, k_d, k_r , KKKS, E2, KKKSE2, KKK) \parallel enz_reaction(k_a, k_d, k_r , KK, KKKS, KKKKS, KKKP) \parallel enz_reaction(k_a, k_d, k_r , KKP, KKP1, KKKS, KKKP) \parallel enz_reaction(k_a, k_d, k_r , KRP, KKP1, KKKS, KKPKNS, KKPP) \parallel enz_reaction(k_a, k_d, k_r , K, KKPP, KKKP, KKPP, KKP) \parallel enz_reaction(k_a, k_d, k_r , K, KKPP, KKPP, KKPP, KPP) \parallel enz_reaction(k_a, k_d, k_r , K, KPP, KKPP, KPPN, KPPN, KPPN) \parallel enz_reaction(k_a, k_d, k_r , KPP, KPPN, KPPNP1, KPPNP1





The gene machine



The instruction set



$$\begin{array}{l} \operatorname{null_gate}(\textit{k}_{\textit{p}}, \textit{X}) : - \\ \operatorname{tell}_{\textit{k}_{\textit{p}}}(\textit{X} = \textit{X} + 1).\operatorname{null_gate}(\textit{k}_{\textit{p}}, \textit{X}) \end{array}$$

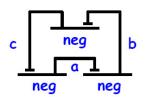
$$\begin{array}{l} \operatorname{pos_gate}(k_{p},k_{e},k_{f},X,Y): - \\ \operatorname{tell}_{k_{p}}(X=X+1).\operatorname{pos_gate}(k_{p},k_{e},k_{f},X,Y) \\ +\operatorname{ask}_{f}(k_{e},\gamma)(\textit{true}).\operatorname{tell}_{k_{B}}(X=X+1).\operatorname{pos_gate}(k_{p},k_{e},k_{f},X,Y) \end{array}$$

$$\begin{aligned} & \text{neg_gate}(k_{p}, k_{i}, k_{d}, X, Y) : -\\ & \text{tell}_{k_{p}}(X = X + 1).\text{neg_gate}(k_{p}, k_{i}, k_{d}, X, Y)\\ & + \text{ask}_{r(k_{i}, Y)}(true).\text{ask}_{k_{d}}(true).\text{neg_gate}(k_{p}, k_{i}, k_{d}, X, Y) \end{aligned}$$

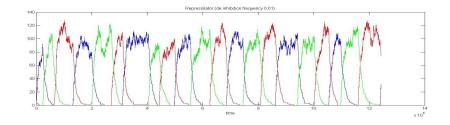
where
$$r(k, Y) = k \cdot Y$$
.

L. Cardelli, A. Phillips, 2005.

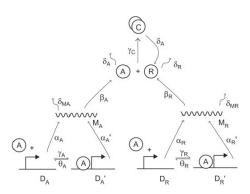
Repressilator



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\begin{array}{l} \text{neg\_gate}(0.1,1,0.0001,A,C) \parallel \\ \text{reaction}(0.0001,[A],[]) \parallel \\ \text{neg\_gate}(0.1,1,0.0001,B,A) \parallel \\ \text{reaction}(0.0001,[B],[]) \parallel \\ \text{neg\_gate}(0.1,1,0.0001,C,B) \parallel \\ \text{reaction}(0.0001,[C],[]) \end{array}
```

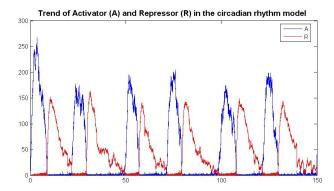


Circadian Clock



Circadian Clock

```
\begin{array}{l} \operatorname{pos\_gate}(\alpha_A,\alpha_A',\gamma_A,\theta_A,M_A,A) \parallel \operatorname{pos\_gate}(\alpha_R,\alpha_R',\gamma_R,\theta_R,M_R,A) \parallel \\ \operatorname{reaction}(\beta_A,[M_A],[A]) \parallel \operatorname{reaction}(\delta_{MA},[M_A],[]) \parallel \\ \operatorname{reaction}(\beta_R,[M_R],[R]) \parallel \operatorname{reaction}(\delta_{MR},[M_R],[R]) \parallel \\ \operatorname{reaction}(\gamma_C,[A,R],[AR]) \parallel \operatorname{reaction}(\delta_A,[AR],[R]) \parallel \\ \operatorname{reaction}(\delta_A,[A],[]) \parallel \operatorname{reaction}(\delta_R,[R],[]) \end{array}
```



Conclusions

- We have introduced a stochastic version of CCP, with functional rates.
- We showed that sCCP may be used for modeling biological systems, defining libraries for biochemical reactions and gene regulatory networks.
- We showed that non-constant rates allow to use more complex chemical kinetics than mass action one.

The End

THANKS FOR THE ATTENTION!

QUESTIONS?