

AUTOMATED REASONING

Agostino Dovier

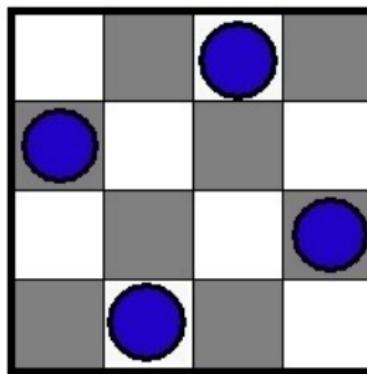
Università di Udine
CLPLAB

Udine, November 2016

Modeling CSP with ASP

N-QUEENS

4-Queens: put 4 queens on a 4×4 chessboard in such a ways as they don't attack each other.



N-QUEENS

Let us define the predicate `queen/2`, where `queen(i,j)` holds iff in the cell (i,j) there is a queen.

Domain predicate (using a constant *n*):

```
numero(1..n).
```

For each column *j* there is exactly one queen:

```
1 {queen(I,J) : numero(I)} 1 :- numero(J).
```

N-QUEENS

HORIZONTAL ATTACK

```
:– numero(I), numero(J1), numero(J2),  
    J1 != J2,  
    queen(I, J1), queen(I, J2) .
```

Read as *It is not possible that there is a row I containing two queens (in column J1 and in column J2)*

In logic:

$$(\exists I)(\exists J_1)(\exists J_2)(J_1 \neq J_2 \wedge \text{queen}(I, J_1) \wedge \text{queen}(I, J_2)) \rightarrow \text{false}$$

Or, equivalently:

$$(\forall I)(\forall J_1)(\forall J_2)((J_1 \neq J_2 \wedge \text{queen}(I, J_1) \wedge \text{queen}(I, J_2)) \rightarrow \text{false})$$

N-QUEENS

HORIZONTAL ATTACK

```
:– numero(I), numero(J1), numero(J2),  
    J1 != J2,  
    queen(I,J1), queen(I,J2).
```

Read as *It is not possible that there is a row I containing two queens (in column J1 and in column J2)*

In logic:

$$(\exists I)(\exists J_1)(\exists J_2) (J_1 \neq J_2 \wedge \text{queen}(I, J_1) \wedge \text{queen}(I, J_2)) \rightarrow \text{false}$$

Or, equivalently:

$$(\forall I)(\forall J_1)(\forall J_2) ((J_1 \neq J_2 \wedge \text{queen}(I, J_1) \wedge \text{queen}(I, J_2)) \rightarrow \text{false})$$

N-QUEENS

DIAGONAL ATTACK

```
:‐ numero(I1), numero(I2), numero(J1), numero(J2),  
    J1 != J2,  
    queen(I1,J1), queen(I2,J2),  
    | I1-I2 | = | J1-J2 |.
```

Read as *It is not possible that there are numbers I_1, I_2, J_1, J_2 with $J_1 \neq J_2$, such that there the queens in (I_1, J_1) and (I_2, J_2) attack each other since $|I_1 - I_2| = |J_1 - J_2|$.*

What happens if we changed $I_1 \neq I_2$ with $I_1 < I_2$?

N-QUEENS

DIAGONAL ATTACK

```
:‐ numero(I1), numero(I2), numero(J1), numero(J2),  
    J1 != J2,  
    queen(I1,J1), queen(I2,J2),  
    | I1-I2 | = | J1-J2 |.
```

Read as *It is not possible that there are numbers I_1, I_2, J_1, J_2 with $J_1 \neq J_2$, such that there are queens in (I_1, J_1) and (I_2, J_2) attack each other since $|I_1 - I_2| = |J_1 - J_2|$.*

What happens if we changed $I_1 \neq I_2$ with $I_1 < I_2$?

OTHER EXAMPLES

Coloring
Schur
Hamiltonian

HAMMING CODES



HAMMING CODES

```
word(1..k).  
position(1..n).  
dmin(d).  
  
%% Values (you can enlarge the alphabet)  
val(0).  
val(1).
```

HAMMING CODES

We wish to define the predicate

tupla(indice, posizione, valore)

```
l{ tupla(I,J,V) : val(V) } l :- word(I), position(J) .  
  
distanza(I1,I2,S) :-  
    word(I1), word(I2), I1 < I2,  
    S = #count{ J : tupla(I1,J,V1), tupla(I2,J,V2), V1 != V2,  
               position(J), val(V1), val(V2)} .  
  
badcode :-  
    word(I1), word(I2), I1 < I2, dmin(D),  
    distanza(I1,I2,S),  
    S < D.  
:- badcode.
```

HAMMING CODES

We wish to define the predicate

tupla(indice, posizione, valore)

```
1{ tupla(I,J,V) : val(V) } 1 :- word(I), position(J).
```

```
distanza(I1,I2,S) :-  
    word(I1), word(I2), I1 < I2,  
    S = #count{ J : tupla(I1,J,V1), tupla(I2,J,V2), V1 != V2,  
                position(J), val(V1), val(V2) }.
```

```
badcode :-  
    word(I1), word(I2), I1 < I2, dmin(D),  
    distanza(I1,I2,S),  
    S < D.  
:- badcode.
```

HAMMING CODES

Symmetry breaking

```
%%% Zero is in the code
```

```
tupla(1,J,0) :- position(J).
```

```
%%% Tuples are lexico ordered
```

```
ordered(I1,I2) :-  
    word(I1), word(I2), I1 < I2,  
    ordered(I1,I2,1).
```

```
ordered(I1,I2,J) :-  
    position(J),  
    word(I1), word(I2), I1 < I2,  
    val(V1), val(V2),  
    tupla(I1,J,V1), tupla(I2,J,V2),  
    V1 < V2.
```

```
ordered(I1,I2,J) :-  
    position(J), position(J+1),  
    word(I1), word(I2), I1 < I2,  
    val(V),  
    tupla(I1,J,V), tupla(I2,J,V),  
    ordered(I1,I2,J+1).
```

```
:-     word(I1), word(I2), I1 < I2, not ordered(I1,I2).
```

HAMMING CODES

Output

```
% Complete with nonsense values if n < 10
extrapos(n+1..10).
tupla(I,J,x) :- word(I), extrapos(J).

mostra(X1,X2,X3,X4,X5,X6,X7,X8,X9,X10) :-
    word(I),
    tupla(I,1,X1),
    tupla(I,2,X2),
    tupla(I,3,X3),
    tupla(I,4,X4),
    tupla(I,5,X5),
    tupla(I,6,X6),
    tupla(I,7,X7),
    tupla(I,8,X8),
    tupla(I,9,X9),
    tupla(I,10,X10).
#show mostra/10.
```

ESERCISE

- Given a chessboard $n \times n$, and two numbers h and k , put h castles (torri) and k bishops (alfieri) in such a way that they do not attack each other.
- Given a chessboard $m \times n$, find a path for a knight (cavalllo) that visits each cell exactly once (starting and arrival places are of course different).
- Consider the Conway's game of life:
https://it.wikipedia.org/wiki/Gioco_della_vita. Consider a board of size $n \times n$ and write an ASP program that finds a stable configuration (see the game rules) with exactly h live cells.