

AUTOMATED REASONING

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WELL-KNOWN.

SEE THE MINIZINC MODELING

Observe the definition of predicates in Minizinc

WELL-KNOWN.
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HAMMING CODES

Given $n, k, d \in \mathbb{N}$, is there a **Code** made of k different n -tuples (of bits) such that each pair of them has a distance of at least d ?

For instance $n = 5, k = 4, d = 3$:

	0	0	0	0	0
	0	0	1	1	1
	1	1	0	1	0
	1	1	1	0	1

SEE THE MINIZING MODELING

Observe the importance of symmetry breaking.

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There are q courses, and each course $i \in \{1, \dots, q\}$ consists of

- k_i lectures, and p periods $1, \dots, p$.
- For all $i \in \{1, \dots, q\}$ all lectures $\ell \in \{1, \dots, k_i\}$ must be assigned to a period $j \in \{1, \dots, p\}$ in such a way that the following constraints are satisfied:
 - 1 Conflicts: There is a conflict matrix M such that $M[i, j] = 1$ if courses i and j have common students. Lectures of courses i and j must be all scheduled at different times
 - 2 Availabilities: There is an availability binary matrix A such that $A[i, j] = 1$ then lectures of course i cannot be scheduled at period j .
 - 3 Rooms: There are r rooms available. At most r lectures can be scheduled at period k , for each $k \in \{1, \dots, p\}$.

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TRAVELING SALESMAN PROBLEM

Input: a complete directed graph with a weight on each edge.

Problem: find the Hamiltonian circuit of minimum cost.

It is crucial to use the global constraint `circuit` that constrains the elements of x to define a circuit where $x[i] = j$ mean that j is the successor of i .

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